

Uncertainty Analysis for Complex Systems: Algorithms for Practical Systems

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Overview

- **Stochastic PDE: Re-formulation**
- **Solution Strategies: generalized polynomial chaos (gPC)**
- **Application: Epistemic uncertainty analysis**

(Re-)Formulation of PDE: Input Parameterization

$$\frac{\partial u}{\partial t}(t, x) = \mathcal{L}(u) \quad + \quad \text{boundary/initial conditions}$$

- **Goal:** To characterize the random inputs by a set of random variables
 - Finite number
 - Mutual independence
- **If inputs == parameters**
 - Identify the (smallest) independent set
 - Prescribe probability distribution
- **Else if inputs == fields/processes**
 - Approximate the field by a function of finite number of RVs
 - Well-studied for Gaussian processes
 - Under-developed for non-Gaussian processes
 - Examples: Karhunen-Loeve expansion, spectral decomposition, etc.

$$a(x, \omega) \approx \mu_a(x) + \sum_{i=1}^d \tilde{a}_i(x) Z_i(\omega)$$

The Reformulation

- **Stochastic PDE:**

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \quad \text{boundary/initial conditions}$$

- **Solution:** $u(t, x, Z) : [0, T] \times \bar{D} \times \mathbb{R}^{n_z} \rightarrow \mathbb{R}$

- Uncertain inputs are characterized by n_z random variables Z

- Probability distribution of Z is prescribed

$$F_Z(s) = \Pr(Z \leq s), \quad s \in \mathbb{R}^{n_z}$$



Non-trivial task

Generalized Polynomial Chaos (gPC)

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Focus on dependence on Z :** $u(\bullet, Z) : \mathbb{R}^{n_z} \rightarrow \mathbb{R}$

- **N^{th} -order gPC expansion:**

$$u_N(t, x, Z) \triangleq \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z), \quad \# \text{ of basis} = \binom{n_z + N}{N}$$

- **Orthogonal basis:** $\int \Phi_i(Z) \Phi_j(Z) \rho(Z) dZ = \delta_{ij}$

- **Basis functions:**

- Hermite polynomials: seminal work by *R. Ghanem*
- General orthogonal polynomials (*Xiu & Karniadakis, 2002*)

- **Properties:**

- Rigorous mathematics
- High accuracy, fast convergence
- Curse-of-dimensionality

- **Numerical Approaches:**

- Galerkin vs. collocation

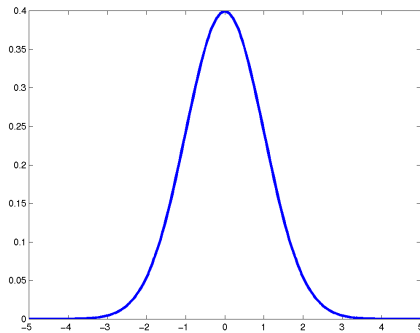
gPC Basis

▪ **Expectation:**

$$\mathbb{E}(g(Z)) = \int_{\mathbb{R}} g(z)\rho(z) dz$$

▪ **Orthogonality:**

$$\int \Phi_i(z)\Phi_j(z)\rho(z) dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$$

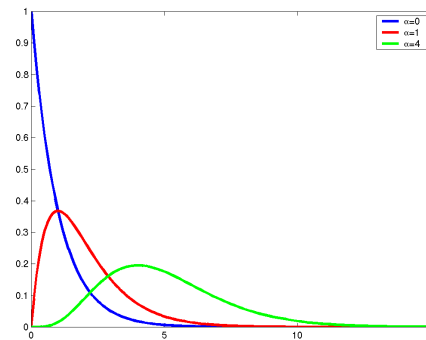


Gaussian distribution

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



Hermite polynomial

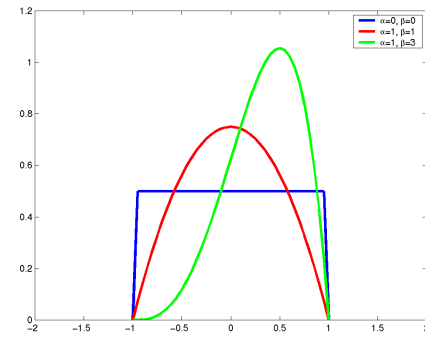


Gamma distribution

$$\int_0^{\infty} \Phi_i(z)\Phi_j(z)e^{-z} dz = \delta_{ij}$$



Laguerre polynomial



Beta distribution

$$\int_{-1}^1 \Phi_i(z)\Phi_j(z) dz = \delta_{ij}$$



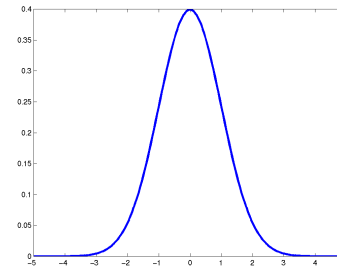
Legendre polynomial

gPC Basis: the Choices

▪ **Orthogonality:** $\int \Phi_i(z)\Phi_j(z)\rho(z)dz = \mathbb{E}[\Phi_i(Z)\Phi_j(Z)] = \delta_{ij}$

▪ **Example: Hermite polynomial**

$$\int_{-\infty}^{\infty} \Phi_i(z)\Phi_j(z)e^{-z^2} dz = \delta_{ij}$$



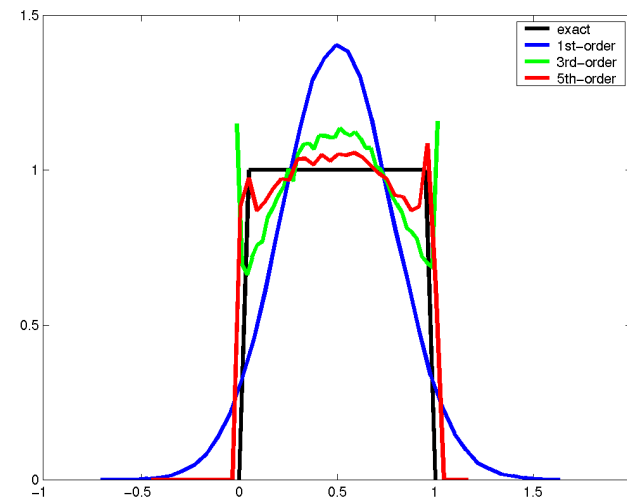
▪ **The polynomials:** $Z \sim N(0,1)$

$$\Phi_0 = 1, \quad \Phi_1 = Z, \quad \Phi_2 = Z^2 - 1, \quad \Phi_3 = Z^3 - 3Z, \quad \dots$$

▪ **Approximation of arbitrary random variable:** Requires L^2 integrability

▪ **Example:** Uniform random variable

- Convergence
- Non-optimal
- First-order Legendre is exact

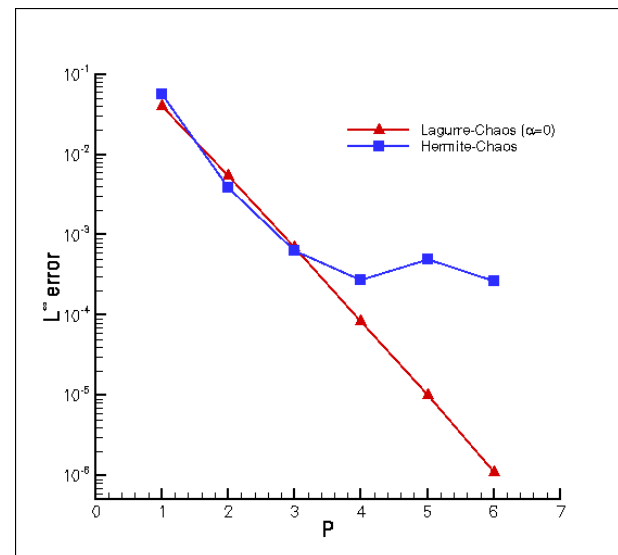
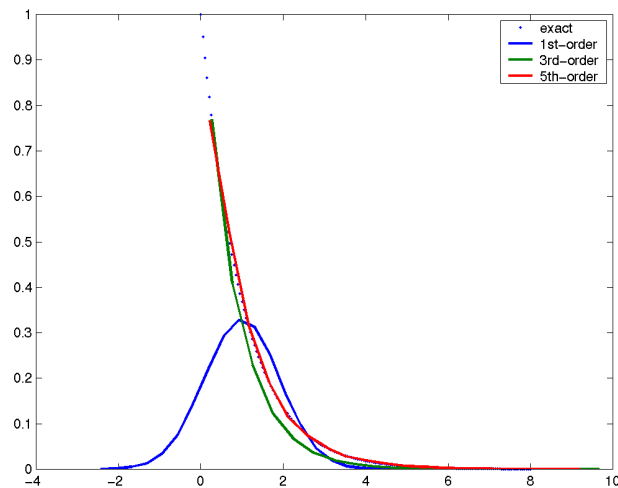


Computational Efficiency

■ First-order ODE: exponential random input

Error	Monte Carlo Method (# of realizations)	Generalized Polynomial Chaos (# of expansion terms)	Speed-up factor
4%	100	1	100
1.1%	1,000	2	500
0.05%	9,800	3	3,267

Effect of non-optimal basis



Stochastic Galerkin

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Galerkin method:** Seek

$$u_N(t, x, Z) \triangleq \sum_{|\mathbf{k}|=0}^N \hat{u}_{\mathbf{k}}(t, x) \Phi_{\mathbf{k}}(Z)$$

Such that

$$\mathbb{E} \left[\frac{\partial u_N}{\partial t}(t, x, Z) \Phi_{\mathbf{m}}(Z) \right] = \mathbb{E} \left[\mathcal{L}(u_N) \Phi_{\mathbf{m}}(Z) \right], \quad \forall |\mathbf{m}| \leq N$$

- **The result:**
 - Residue is orthogonal to the gPC space
 - A set of deterministic equations for the coefficients
 - The equations are usually coupled – requires new solver

Stochastic Collocation

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) + \text{boundary/initial conditions}$$

- **Collocation:** To satisfy governing equations at selected nodes
 - Allow one to use existing deterministic codes repetitively

- **Sampling:** (solution statistics only)
 - Random (Monte Carlo)
 - Deterministic (lattice rule, tensor grid, cubature)

- **Stochastic collocation:** To construct **polynomial approximations**
 - Node selection is critical to efficiency and accuracy
 - More than sampling

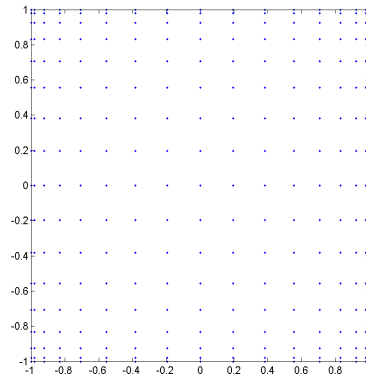
Stochastic Collocation: Interpolation

$$\frac{\partial u}{\partial t}(t, x, Z) = \mathcal{L}(u) \quad + \text{ boundary/initial conditions}$$

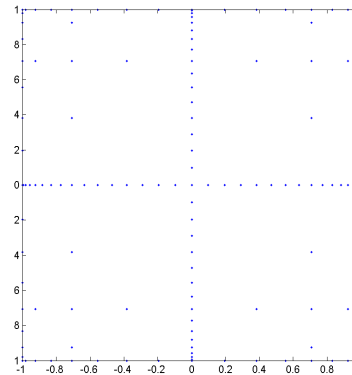
- **Definition:** Given a set of nodes and solution ensemble, find u_N in a proper polynomial space, such that $u_N \approx u$ in a proper sense.

- **Interpolation Approaches:**
$$u_N(Z) = \sum_{j=1}^Q u(Z^j) L_j(Z)$$
$$L_i(Z^j) = \delta_{ij}, \quad 1 \leq i, j \leq Q$$

- Dimension-by-dimension space filling



Tensor grids: inefficient



Sparse grids: more efficient

(Xiu & Hesthaven, SIAM J. Sci. Comput., 05)

Stochastic Computation: The Landscape

- **Realistic Large-scale Complex Systems:**

- Complex physics → highly nonlinear systems
- Large number of random variables
- (Extremely) time consuming simulations
- Legacy codes (nearly impossible to re-write)

- **Stochastic Galerkin:**

- Difficult to implement
- Good mathematical properties

- **Stochastic collocation is more proper:**

- Easy to implement → virtually no coding effort
- Nonlinearity poses no additional difficulties

Epistemic Uncertainty: Setup

• **Governing Equation:**

$$\left\{ \begin{array}{ll} \frac{\partial v}{\partial t}(t, x, Z) = \mathcal{L}(v), & D \times (0, T] \times I_Z \\ \mathcal{B}(v) = 0, & \partial D \times [0, T] \times I_Z \\ v = v_0, & D \times \{t = 0\} \times I_Z \end{array} \right.$$

$$v(x, t, Z) : \bar{D} \times [0, T] \times I_Z \rightarrow \mathbb{R} \quad I_Z \subseteq \mathbb{R}^d$$

- **Epistemic uncertainty:**
 - Distribution of Z is not fully known
- **“Some” prior knowledge:**

$$I_{Z_i} = [\alpha_i, \beta_i], \quad -\infty \leq \alpha_i < \beta_i \leq \infty$$

$$I_Z \subseteq \times_{i=1}^d I_{Z_i}$$

- **Remark:** Z_i can be dependent, unbounded, and I_Z can be much smaller.

Encapsulation

- **Goal:** To “encapsulate” each variable

- For each Z_i (potentially unbounded), find a **bounded interval** to “capture” it.
- Requires modeling effort

- **Overwhelming probability condition:**

For each $I_{Z_i} = [\alpha_i, \beta_i]$, $\alpha_i < \beta_i$, find a bounded interval

$$I_{X_i} = [a_i, b_i], \quad -\infty < a_i < b_i < \infty,$$

such that

$$\Pr(Z_i \in I_i^-) \leq \delta_i,$$

where $\delta_i \geq 0$, and I_i^- is the difference set

$$I_i^- = I_{Z_i} \Delta I_{X_i} = (I_{Z_i} \cup I_{X_i}) \setminus (I_{Z_i} \cap I_{X_i})$$

- If Z_i is bounded, it is “easier” to do
- If Z_i is unbounded, X_i needs to be “big” enough

Encapsulation (cont'd)

- For each variable: $I_{Z_i} = [\alpha_i, \beta_i], \quad -\infty \leq \alpha_i < \beta_i \leq \infty$

$$I_{X_i} = [a_i, b_i], \quad -\infty < a_i < b_i < \infty$$

$$\Pr(Z_i \in I_i^-) \leq \delta_i$$

- For all variables: $I_Z \subseteq \times_{i=1}^d I_{Z_i}$

$$I_X = \times_{i=1}^d I_{X_i} = \times_{i=1}^d [a_i, b_i]$$

$$I^+ = I_Z \cup I_X, \quad I^o = I_Z \cap I_X$$

$$I^- = I_Z \Delta I_X = I^+ \setminus I^o \text{ (difference set)}$$

- Overwhelming probability condition:

$$\Pr(Z_i \in I^-) \leq \delta, \quad \delta = 1 - (1 - \delta_i)^d$$

- **Reminder:** I_X may not overlap I_Z

Encapsulation Problem

• **Original Problem:**

$$\left\{ \begin{array}{ll} \frac{\partial v}{\partial t}(t, x, Z) = \mathcal{L}(v), & D \times (0, T] \times I_Z \\ \mathcal{B}(v) = 0, & \partial D \times [0, T] \times I_Z \\ v = v_0, & D \times \{t = 0\} \times I_Z \end{array} \right.$$

• **Encapsulation Problem:**

$$\left\{ \begin{array}{ll} \frac{\partial u}{\partial t}(t, x, X) = \mathcal{L}(u), & D \times (0, T] \times I_X \\ \mathcal{B}(u) = 0, & \partial D \times [0, T] \times I_X \\ u = v_0, & D \times \{t = 0\} \times I_X \end{array} \right.$$

○ Solution in a hypercube: $u(x, t, X) : \bar{D} \times [0, T] \times I_X \rightarrow \mathbb{R}$

$$I_X = [a_i, b_i]^d (= [-1, 1]^d, = [0, 1]^d)$$

• **Assumption:**

$$u(\cdot, \xi) = v(\cdot, \xi), \quad \forall \xi \in I^o$$

Solution Strategy of the Encapsulation Problem

- **Encapsulation Problem:**

$$\left\{ \begin{array}{ll} \frac{\partial u}{\partial t}(t, x, X) = \mathcal{L}(u), & D \times (0, T] \times I_X \\ \mathcal{B}(u) = 0, & \partial D \times [0, T] \times I_X \\ u = v_0, & D \times \{t = 0\} \times I_X \end{array} \right.$$

- **Solution strategy:** Controllability on point-wise error

$$\varepsilon_n = \|u - u_n\|_{L^\infty(I_X)} \rightarrow 0, \quad n \rightarrow \infty$$

- u_n is a good approximation in the entire domain I_X (hypercube)
 - Can “sample” u_n accurately for all realizations
 - No probability distribution is assigned in I_X .
 - **Convergence is a mathematical preference, not a practical necessity**
- Requirement on error control is strong but achievable
 - Sparse grid collocation (with sufficient regularity)
 - Polynomial Galerkin (e.g., Chebyshev) methods are possible
 - Without sufficient regularity --- multi-element approach

Solution “Statistics”

- Solution of the original problem:

$$v(\cdot, Z) : I_Z \rightarrow \mathbb{R} \qquad \mu = \int_{I_Z} v(s) \rho_Z(s) ds$$

- Solution in the hypercube:

$$I^o = I_Z \cap I_X$$
$$u_n(\cdot, X) : I_X \rightarrow \mathbb{R} \qquad \mu_n = \int_{I^o} u_n(s) \rho_Z(s) ds$$

Theorem : Assume $v(Z)$ is bounded and let $C_v = \|v\|_{L^\infty(I_Z)}$. Let u_n be an approximation to the solution of the encapsulation problem $u(X)$, s.t.,

$$\varepsilon_n = \|u - u_n\|_{L^\infty(I_X)}.$$

Then the approximation of the mean solution satisfies

$$|\mu - \mu_n| \leq \varepsilon_n + C_v \cdot \delta$$

Numerical Example

- **Original Problem:**

$$\frac{d^2 v}{dt^2}(t, Z) + \gamma \frac{dv}{dt} + kv = f \cos(\omega t), \quad v(0) = v_0, \quad \frac{dv}{dt}(0) = v_1$$

$$Z = (\gamma, k, f, \omega, v_0, v_1) \in \mathbb{R}^6$$

- **Encapsulation Problem:**

$$\frac{d^2 u}{dt^2}(t, X) + X_1 \frac{du}{dt} + X_2 u = X_3 \cos(X_4 t), \quad u(0) = X_5, \quad \frac{du}{dt}(0) = X_6$$

$$X = (X_1, \dots, X_6) \in [-1, 1]^6$$

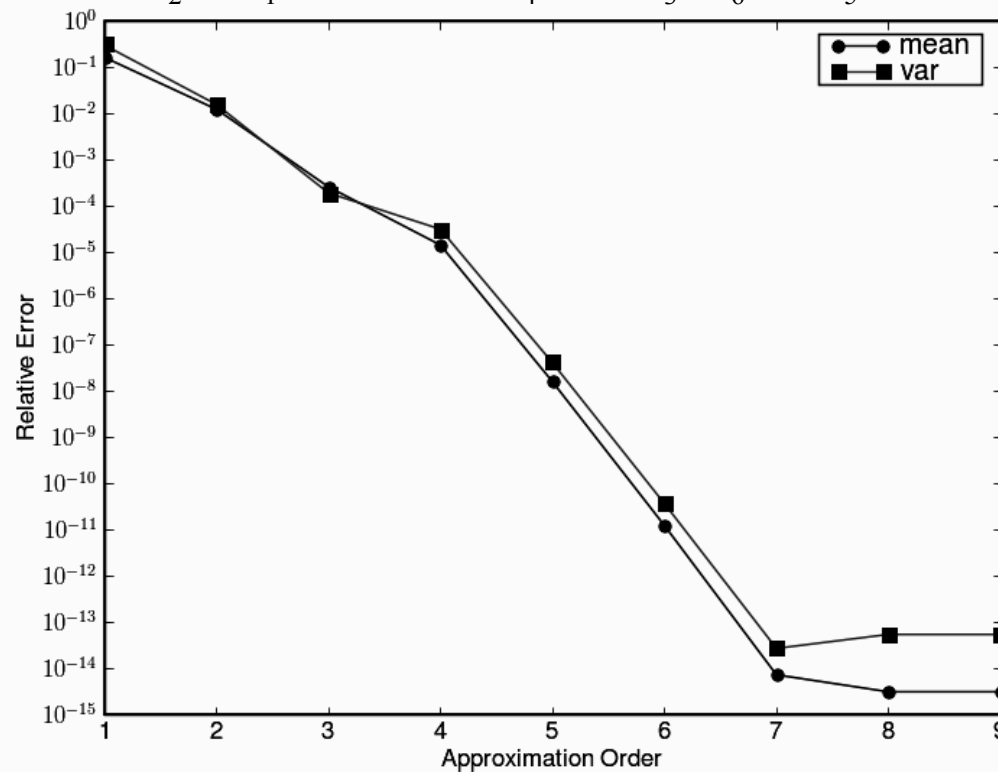
- Solved by 6-dimensional sparse grid collocation for $t=20$

Dependent Inputs

$$\frac{d^2 u}{dt^2}(t, Z) + Z_1 \frac{du}{dt} + Z_2 u = Z_3 \cos(Z_4 t), \quad u(0) = Z_5, \quad \frac{du}{dt}(0) = Z_6$$

$Z_1 \sim \text{beta}(0.08, 0.12, 3, 2)$, $Z_3 \sim \text{beta}(0.08, 0.1, 1, 1)$, $Z_5 \sim \text{uniform}(0.45, 0.55)$, independent

$$Z_2 = Z_1^2 / 4 + 0.01, \quad Z_4 = 10Z_3, \quad Z_6 = (Z_5 - 0.5)$$



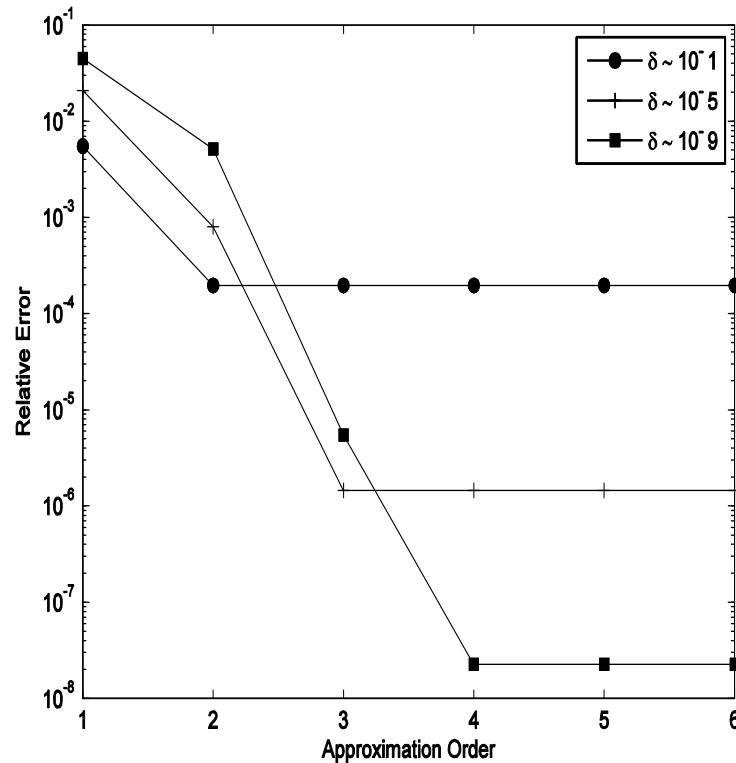
- 4d problem is solved in 6d, without “cutoff”

Unbounded Inputs: Effect of “Cutoff”

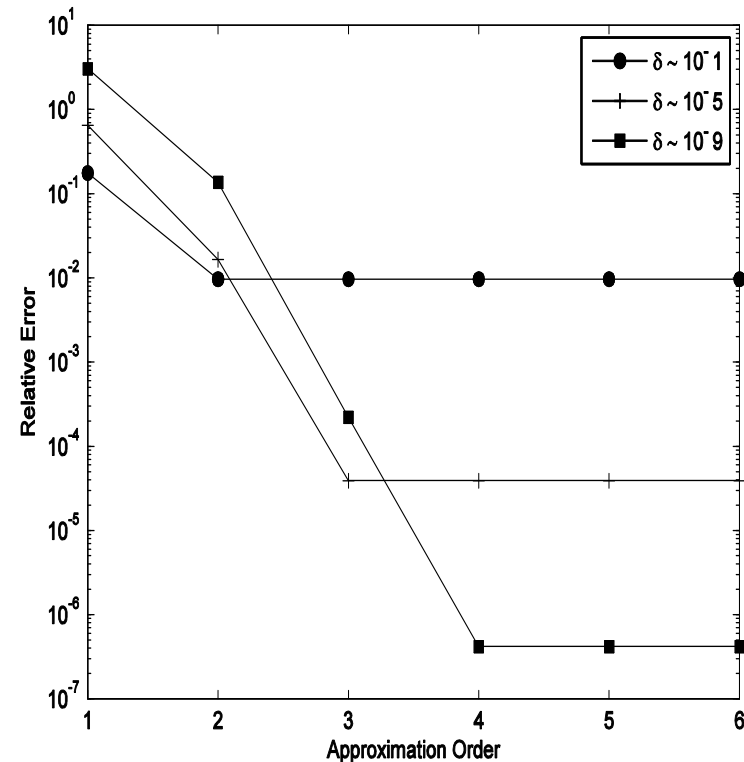
$$\frac{d^2u}{dt^2}(t,Z) + Z_1 \frac{du}{dt} + Z_2 u = Z_3 \cos(Z_4 t), \quad u(0) = Z_5, \quad \frac{du}{dt}(0) = Z_6$$

Gaussian: $Z \sim \mathcal{N}(\mathbf{0}, \mathbf{C})$, $\mathbf{C} \in \mathbb{R}^{6 \times 6}$ is the covariance matrix

$I_Z = \mathbb{R}^6$, $I_X = [-a, a]^6$, then $\delta > 0$



Mean

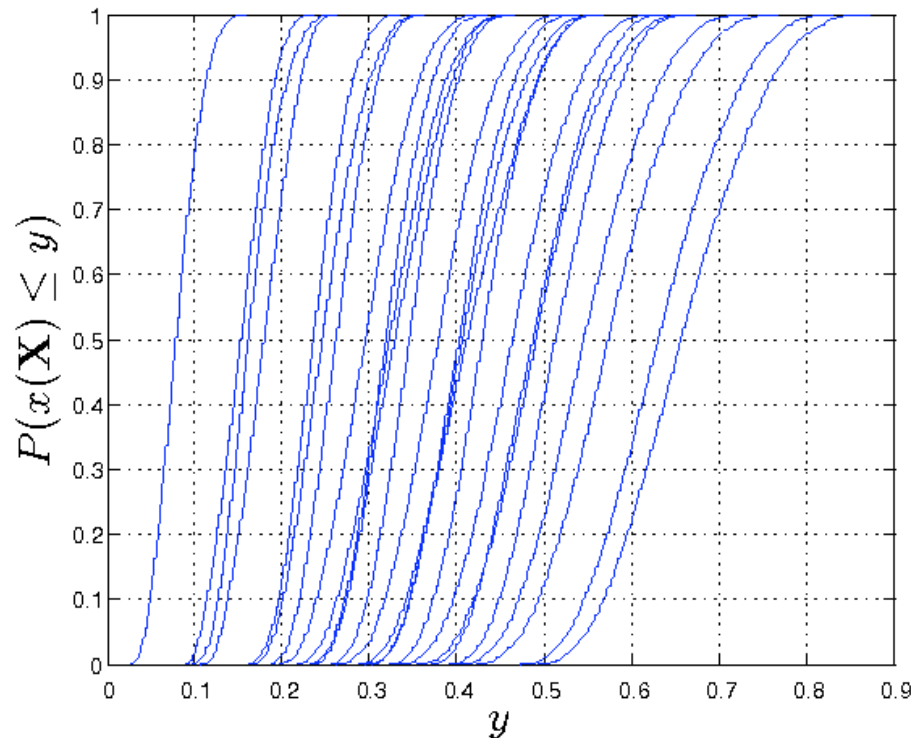


Variance

Mixed Aleatory and Epistemic Case

$$\frac{d^2u}{dt^2}(t,Z) + Z_1 \frac{du}{dt} + Z_2 u = Z_3 \cos(Z_4 t), \quad u(0) = Z_5, \quad \frac{du}{dt}(0) = Z_6$$

- Aleatory: $Z_1 \sim \text{beta}(0,1,0,0)$, $Z_2 = Z_1^2 / 4 + 0.01$, $Z_3 \sim \text{beta}(0,1,1,1)$, $Z_5 \sim \text{beta}(0,1,2,1)$
- Epistemic: $Z_4 \in [0.8, 1.2]$, $Z_6 \in [-0.05, 0.05]$



CDF of the aleatory RVs at 25 prescribed values of the epistemic variables

- Solution obtained by sparse grid collocation via simultaneous construction (5-d)

Epistemic Uncertainty: Think “Outside the Box”

- **Original Problem:**

$$\left\{ \begin{array}{ll} \frac{\partial v}{\partial t}(t, x, Z) = \mathcal{L}(v), & D \times (0, T] \times I_Z \\ \mathcal{B}(v) = 0, & \partial D \times [0, T] \times I_Z \\ v = v_0, & D \times \{t = 0\} \times I_Z \end{array} \right.$$

- **Encapsulation Problem:**

$$\left\{ \begin{array}{ll} \frac{\partial u}{\partial t}(t, x, X) = \mathcal{L}(u), & D \times (0, T] \times I_X \\ \mathcal{B}(u) = 0, & \partial D \times [0, T] \times I_X \\ u = v_0, & D \times \{t = 0\} \times I_X \end{array} \right.$$

- **Question: Does I_X have to be a hyber-box?**

- ✓ **No. I_X can be unbounded too.**
- ✓ **Numerical solution can converge in L^p norm. (More practical)**
- ✓ **Additional constraints on the measures are needed.**

Reference:

- J. Jakeman, M. Eldred, D. Xiu, “Numerical Approach for Quantification of Epistemic Uncertainty”, *Journal of Computational Physics*, vol. 229, pp. 4648-4663, 2010.
- X. Chen, E.-J. Park, D. Xiu, Preprint, 2011.

Summary

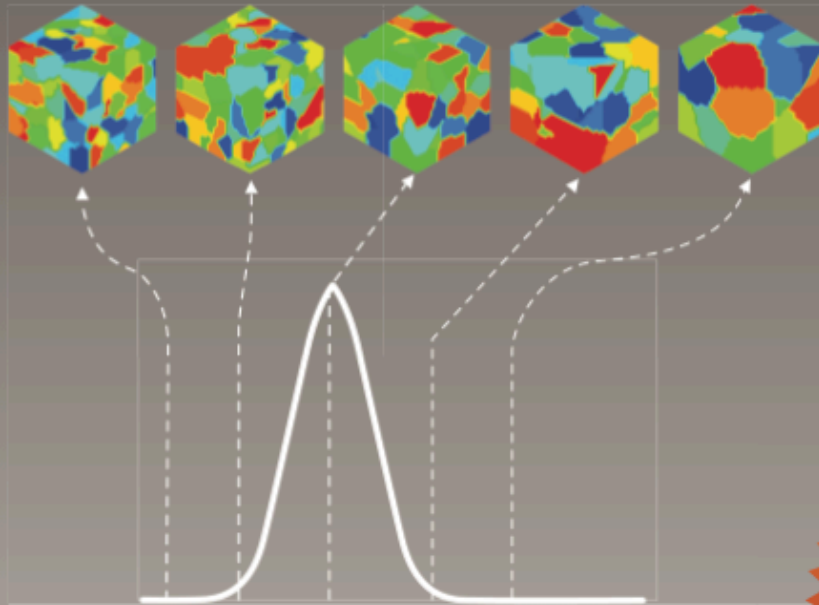
- **Uncertainty Analysis:** To provide improved prediction
 - Input characterization
 - Uncertainty propagation
 - Post processing
- **Generalized polynomial chaos (gPC)**
 - Multivariate approximation theory
- **Important directions:**
 - Approximation theory in HIGH dimensions
 - Combination with data
- **Data, *any* data, can help**

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